ZIPPER GRIDS FOR FORCE AND MOMENT COMPUTATION ON OVERSET GRIDS

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Abstract

A scheme is developed to account for the overlapped zones on overset structured surface grids when computing forces and moments for the surface. Grid points in the overlapped zone between grids are blanked. The resulting gaps are then filled by 'zipper grids' which are strings of triangular cells that conform to the original surface. A hybrid composite surface grid is formed which consists of non-overlapping quadrilaterals and triangles, where the quadrilaterals are the unblanked cells in the original surface grids and the triangles are the newly generated triangular cells in the gaps. Pressure and viscous stresses which are known at the grid points can then be integrated accurately on the hybrid surface. Results obtained by applying the current scheme to numerical solutions are compared with data from wind tunnel tests for several complex configurations. These include the Space Shuttle Launch Vehicle, a subsonic transport, and a supersonic oblique all-wing configuration.

1. Introduction

Chimera overset grid methods[1] have been successfully applied in the numerical flow simulations of complex configurations such as the Space Shuttle Launch Vehicle,[2,3,4] the V-22 tiltrotor,[5] the SOFIA Airborne Observatory,[6] and many others. Very good agreement of many flow characteristics has been obtained between the numerical solution and wind tunnel/flight tests. One of the most important set of quantities sought from numerical solutions is the set of force and moment coefficients for the solid surfaces in the flow field. Accurate computation of these quantities is crucial for study of aerodynamics, design optimization, and moving body problems.

Forces and moments are usually computed by integrating the pressure and viscous stresses on the surface grids. Since the surface grids for overset grids are allowed to overlap arbitrarily, the regions in the overlapped zones are counted more than once if integration is performed over all the surface grids. In order for overset grid methods to be widely accepted for use in practical applications, a scheme must be available to accurately extract force and moment information from the numerical solutions. A procedure developed to solve this problem is described in this paper.

The method for computing force and moment coefficients is broken into two main steps performed by two separate computer codes. The first step is to construct a hybrid composite surface grid consisting of non-overlapping quadrilaterals (quads) and triangles (the MIXSUR code), while the second step is to integrate the pressure and viscous stresses on the hybrid grid (the OVERINT code). Generation of the hybrid composite surface grid only needs to be performed once before the flow solver is called. Integration of pressure and viscous stresses can be performed just once at the end of the run to compute the final force and moment coefficients. It can also be performed every N number of iterations to supply the force and moment coefficients as indicators of convergence, or as input to a six-degrees-of-freedom rigid body dynamics package for moving body applications. In both cases, the connections between the quads and the triangles in the hybrid composite surface grid do not change with iterations.

Steps involved in the generation of the hybrid grid are described in §2. The integration scheme is presented in §3. In §4, the results of applying the current scheme on numerical solutions are compared with data obtained in wind tunnel tests, and results from using an alternative force/moment computation scheme.[4] Summary and conclusions are given in §5.

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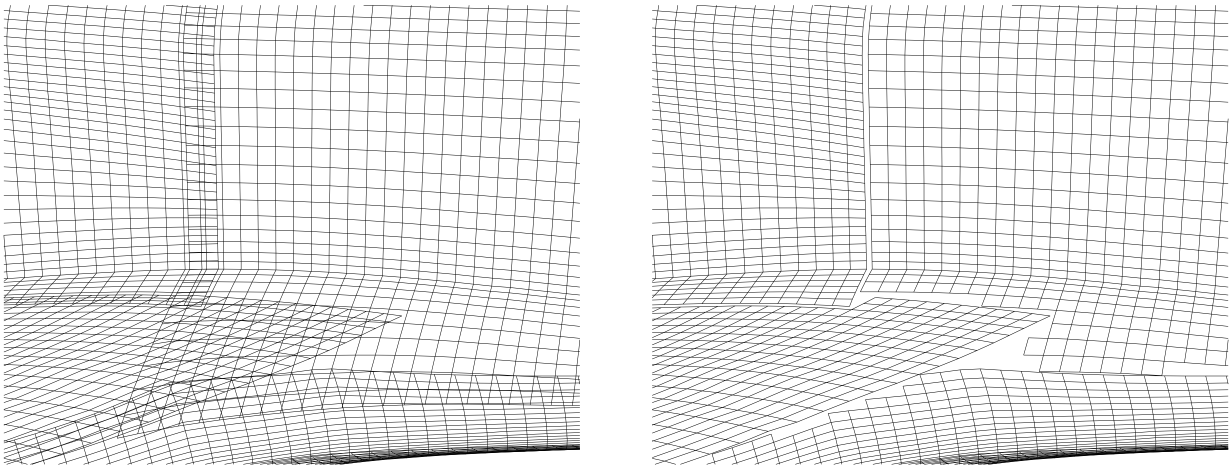


(a) (b)

Figure 1. (a) Surface grids for four overlapping subsets. (b) Surface grids after blanking of points in overlapped regions.

2. Hybrid Composite Grid Generation

An overset grid system consists of a collection of overlapping volume grids together with information on inter-grid connectivity. Let J, K and L be the grid indices of a typical volume grid. The surface domain on which forces and moments are to be computed can be described by a collection of surface subsets. Each surface subset is a subset of a J, K or L = constant slice of a volume grid in the system. For the rest of this paper, a surface subset will be referred to as a subset for short. An example with four overlapping subsets is shown in Figure 1a. A hybrid composite grid will be generated from the collection of subsets. This procedure consists of the following three main automated steps which are described in more detail in the subsections below.

(a) Points in the overlapped regions between subsets are blanked resulting in gaps between subsets. For any overlapping pair of subsets, only points from the coarser of the two subsets are blanked.

(b) Points on the boundaries of the gaps are identified and ordered. This creates a collection of strings of gap boundary points.

(c) Matching pairs of strings of gap boundary points are connected to form a zipper grid which is a single layer of triangular cells that fills the gap between neighboring subsets and conforms to the original surface. In regions where more than two subsets overlap, multiple zipper grids are allowed to converge in a pocket of triangular cells.

An alternative to the above procedure is to retain all the points in all the subsets including those in the overlapped zones and then perform triangulation over the entire surface domain.[7] We have chosen not to do this for the following reason. If there are flow gradients not fully resolved by neighboring grids in the overlapped region, the solution will differ slightly. Use of points from all neighboring grids in this region would result in integration of an unsmooth pressure and/or viscous stress field. Although the hybrid grid generation procedure is more complex in the zipper grid approach, the effect of any solution mismatches is minimized by keeping as many non-overlapping quads as possible and introducing just a single layer of triangular cells. For conservation purposes, hybrid grids have also been used to replace overlapping structured grids in a 2-D or 3-D domain.[8] In that approach, all points in the overlapped region are blanked and the entire overlapped region is triangulated.

2.1. *Grid Point Blanking Scheme in Overlapped Zone*

The integration of forces and moments is performed on the set of non-overlapping quads and triangular cells. The set of non-overlapping quads is derived from the defined subsets where some of the points are blanked to ensure that the quads do not overlap. This information is stored in an integration iblank variable $B_{j,k}$ for each point on the subsets, where $B_{j,k}$ is defined by

$$B_{j,k} = \begin{cases} 1, & \text{if point is used for integration;} \\ 0, & \text{if point is not used for integration,} \end{cases} \tag{2.1}$$

and j, k are grid indices for a subset. A quad is integrated if all four of its vertices have $B_{j,k} = 1$.

Initially, $B_{j,k}$ is set to the value 1 for all points on the subsets. Any point that is a hole point as a result of Chimera inter-grid connectivity[1] will have $B_{j,k}$ set to 0. Points in the overlapped region between subsets are then identified. For any pair of subsets that overlap, grid points on the coarser of the two subsets are blanked, i.e. $B_{j,k}$ set to 0. The coarser subset is defined to be the one with the fewer grid points in the overlapped region.

Points are identified to be inside an overlapped zone by a 'minmax' box approach.[9] A minmax box is a collection of 6 numbers that indicates the smallest and largest x, y and z coordinates of the objects enclosed. A big minmax box is defined for each subset and small minmax boxes are defined for each quad of each subset. Then the blanking procedure goes as follows.

(1) Let S_i and S_j be any two subsets. Given a point in S_i, the big minmax box of S_j is used to quickly determine whether the point falls inside of S_j. If it does, then the small minmax boxes of S_j are used to determine whether the point falls inside any quad in S_j. If such a quad is found, then the point is counted as falling in the overlapped zone between S_i and S_j.

(2) The number of points in S_i that falls in the overlapped zone with S_j is computed and stored in $N(S_i, S_j)$. Similarly, $N(S_j, S_i)$ is also computed.

(3) If $N(S_i, S_j) \leq N(S_j, S_i)$, then all the points in S_i that fall in the overlapped zone with S_j are blanked. This guarantees that all unblanked points in S_i are not contained in any quad in S_j. However, the reverse is not necessarily true. Hence, a further test is necessary to blank any unblanked points of S_j that fall inside an unblanked quad in S_i. The procedure follows similarly with the role of S_i and S_j reversed if $N(S_j, S_i) < N(S_i, S_j)$.

(4) This is performed for all pairs of subsets.

(5) After all the blanking is completed, some loose strings of unblanked points may remain which are surrounded by blanked points. A final test is then performed to blank out these loose strings of points.

Figure 1b shows the remaining non-overlapping quads after the blanking step is performed on the example in Figure 1a. Since the surface grid lines are generally not aligned with the Cartesian coordinate axes, the use of minmax boxes will tend to blank out more points than is necessary. A more accurate but also more expensive test could be used to reduce the size of the gaps created.

2.2. *Stringing of Gap Boundary Points*

The second step in the hybrid grid generation procedure is to identify and order the points on the gap boundaries created by the blanking step. These points fall under two types: (I) points on the boundary of holes defined by the integration iblank variable, and (II) points on the boundary of subsets that lie in the overlapped region with another subset.

Type I points are identified by examining $B_{j,k}$. First, it must be an unblanked point. For an interior point (j, k) to qualify as a Type I point, at least 3 and no more than 7 of its 8 neighboring points must be unblanked. The neighboring points are located at $(j-1, k-1)$, $(j-1, k)$, $(j+1, k)$, $(j, k-1)$, $(j, k+1)$, $(j-1, k+1)$, $(j, k+1)$ and $(j+1, k+1)$. For a point on the boundary or a corner of a subset, it is a Type I point if any one of its 3 or 2 neighbors are blanked, respectively. Type II points are identified by testing if any boundary points on a given subset lie inside any small minmax box of a cell from another subset.

After identifying the gap boundary points, they need to be connected to form ordered strings of points. Two gap boundary points are connected if an edge of a subset exists between them. The connection of the gap boundary points then results in a collection of gap boundary strings (see Figure 1c).

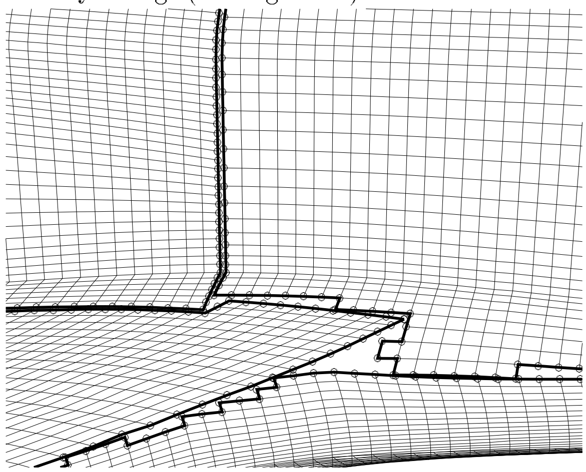
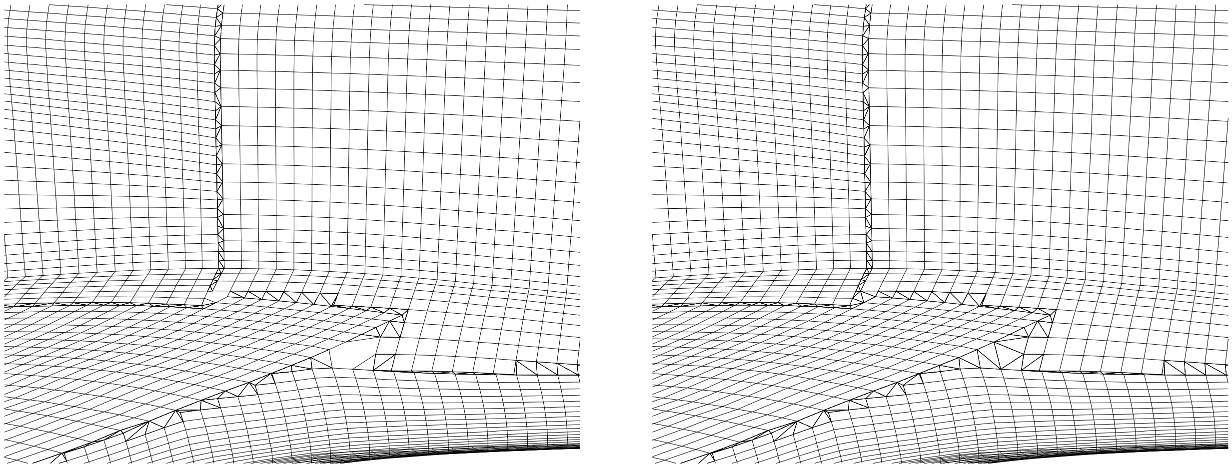


Figure 1. (c) Points on gap boundary strings are indicated by circles joined by thick lines.



(d) (e)

Figure 1. (d) Gaps between pairs of boundary segments are filled by zipper grids where each consists of a single layer of triangular cells. (e) Final hybrid composite surface grid with remaining polygonal gaps filled with triangles.

2.3. *Construction of Zipper Grids*

Triangles are constructed by connecting points between pairs of gap boundary strings resulting in a set of zipper grids. Each zipper grid is a single layer of triangular cells. No new points are introduced in the generation of these cells which are required to conform to the original surface geometry. The steps involved in this procedure are described below.

2.3.1. *Boundary Segment Matching*

A zipper grid is formed by connecting points between segments of two boundary strings. Each boundary string R_i is divided into segments as follows. For each point P_i on R_i, the nearest point on another boundary string R_j is located. The point P_i is now mapped to the boundary string R_j. A segment of R_i is a set of contiguous points on R_i that is mapped to R_j. A segment from R_i that is mapped to R_j is matched with a segment from R_j that is mapped to R_i.

2.3.2. *Zipping*

The matching step identifies segment pairs that should be connected via a zipper grid. Let A and B be a segment pair to be connected. The algorithm used to construct a zipper grid is as follows.

(a) *Selection of initial points*

The first connection between A and B is formed by searching for a starting point in A and a starting point in B to connect to each other. For a non-periodic segment pair A and B, the two end points of segment A and B that are closest to each other are chosen as the starting points of the zipping process. For a periodic segment pair A and B, any two closest points between the two segments can be used as starting points.

(b) *Connection to form triangles*

Let P_i be the ith point on segment A and P_j be the jth point on segment B where $i = j = 1$ at the starting points. Also let $i_{\max}$, and $j_{\max}$ be the number of points on A and B, respectively. For $i < i_{\max}$ and $j < j_{\max}$, P_i may be connected to P_{j+1} or P_j may be connected to P_{i+1}. The first step is to test if the quad formed by the points P_i, P_{i+1}, P_{j+1} and P_j is convex. The connection of P_i to P_{j+1} splits the quad into two triangles. If the normals of these triangles have the same sign, then the connection is valid. A similar test is performed for the P_j to P_{i+1} connection. In cases where only one of the two connections is valid, the valid connection is chosen. In cases where both connections are valid, the quad is convex and the connection with the shortest distance is chosen. However, if both connections are invalid, the quad is crossed inside out. The shortest distance test is still performed and can result in good triangulations in some situations. When bad triangulations are obtained, we will proceed to the next connection but the bad triangles are flagged. When $i = i_{\max}$ is reached (last point of segment A), all remaining points in segment B are connected to the last point in segment A, and vice versa.

(c) *Reconnection to remove bad triangles*

When there is a large discrepancy between the grid resolutions of the neighboring grids in the overlapped zone, and when the grid lines are not aligned, the zipping procedure described above can sometimes produce bad triangles. These triangles are bad in the sense that they overlap some existing quad or good triangle. From experience, typically only less than 1% of the triangles created are bad under the conditions stated above. Some of these bad triangles can be removed and good triangles can be formed by performing the zipping step (b) in the reverse direction in the region of bad triangles. Zipper grids for the example started in §2.1 are shown in Figure 1d.

2.3.3. *Closing of Pockets*

After connecting pairs of gap boundary segments, most of the gaps created by the blanking procedure described in §2.1 are now closed by triangular cells. However, there may be small gaps left in regions where three or more subsets overlap. Each small gap is in the form of a closed polygon where each edge of the boundary is either an unconnected edge of a gap boundary segment or an end edge of one of the zipper grids. These edges are identified and connected to form a closed loop. The gap inside the closed loop is then triangulated without introducing new points using a similar zipping algorithm as described in the previous subsection. First, a starting edge has to be chosen. This is selected to be the edge whose end points contain interior angles that are closest to 90 degress. Zipping then proceeds from the starting edge until the entire polygon is filled with triangles. The complete hybrid composite grid for the example started in §2.1 is shown in Figure 1e.

3. Pressure and Viscous Stress Integration

Integration of the pressure and viscous stresses is performed on the hybrid composite surface consisting of a collection of non-overlapping quadrilaterals and triangles. For each quadrilateral or triangle, the value of the function being integrated is known at the vertices. A number of quadrature rules can be used to perform the integration. For simplicity, the function is assumed to be a constant inside the area of the quadrilateral or triangle where the constant is the arithmetic average of the values of the function at the vertices. This can be written as

$$\int_{polygon} f \, d\vec{A} \approx \left(\frac{1}{N_v} \sum_{i=1}^{Nv} f_i \right) \cdot \vec{n} \Delta A, \qquad (3.1)$$

where f is a scalar function to be integrated, *polygon* is either a quadrilateral or a triangle, N_v is the number of vertices of the polygon, f_i is the value of the scalar function at the vertices, $\vec{n}$ is the unit normal of the polygon, and ΔA is the area of the polygon. A similar expression can be written for the integral of a vector function $\vec{f}$. For a triangle, the normal is uniquely defined. For a quad, the normal is computed by the cross product of the diagonals. The integral in (3.1) is performed over all the unblanked quads and triangles and the results are summed and normalized to obtain the force and moment coefficients.

There are three possible ways to perform the quadrature over a quad by treating it as (1) a 4-point polygon, (2) two triangles by connecting one diagonal of the quad, (3) two triangles by connecting the other diagonal of the quad. Let I_q, I_{t1} and I_{t2} be the value of the integral computed using options (1), (2) and (3), respectively. It can be shown analytically that if the quadrature rule given by (3.1) is used on a parallelogram, then

$$I_{t1} \leq I_q \leq I_{t2}. \qquad (3.1)$$

Hence, for a quad that is a parallelogram, option (1) is the better choice since options (2) and (3) tend to produce undesirable biases. However, as a quad degenerates to a triangle as in the case of two adjacent points becoming very close together, option (2) or (3) is more attractive since the two close points are not counted twice. On typical surface grids for CFD applications, most of the quads are approximately parallelograms in shape. We have therefore chosen option (1) for the quadrature on quads. Cells adjacent to a singular axis point on a surface subset are integrated as triangles.