



LOKI97 - AES1

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Introduction

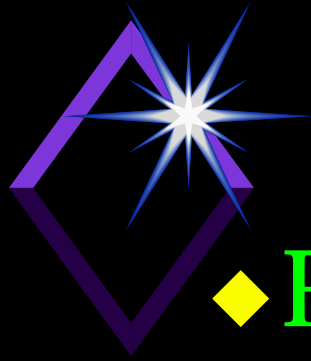
◆ LOKI97

- ◆ new 128-bit private key block cipher
- ◆ based on earlier LOKI89 and LOKI91
- ◆ traditional Feistel S-P design
- ◆ AES candidate



Previous Work - LOKI89

- ◆ 64-bit private key block cipher
- ◆ Brown, Pieprzyk, Seberry 1989/90
 - ◆ see Brown PhD thesis
- ◆ Biham and Shamir, Knudsen
 - ◆ differential cryptanalysis of reduced rounds
 - ◆ full version secure



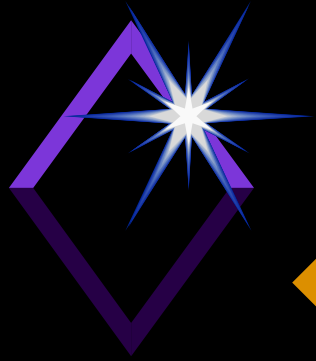
Previous Work - LOKI91

- ◆ Redesigned in 1991
 - ◆ by Brown, Kwan, Pieprzyk, Seberry
 - ◆ better round function and key schedule



Previous Work - LOKI91

- ◆ Knudsen, Biham, Tokita et al
 - ◆ differential cryptanalysis of reduced rounds
 - ◆ linear cryptanalysis of reduced rounds
 - ◆ full version secure from DC *and* LC
 - ◆ some key schedule weaknesses
 - ◆ effective key size approx 60 bits



Design Considerations

- ◆ Motivated by Knudsen 93
 - ◆ no simple relations
(vis data and key)
 - ◆ all keys are equally good
 - ◆ resistant to differential attacks
 - ◆ resistant to linear attacks



(Further) Design Considerations

- ◆ non-linear key schedule
- ◆ highly non-linear round function
- ◆ efficient implementation with tables

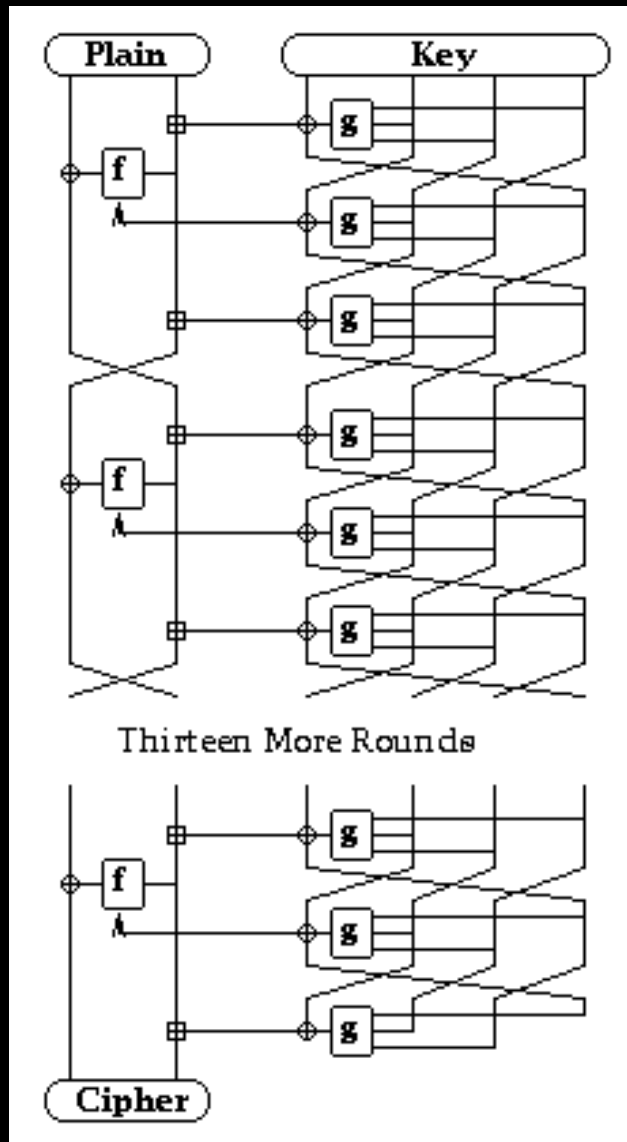


LOKI97 Overview

- ◆ LOKI97 is
 - ◆ private key Feistel S-P block cipher
 - ◆ 128-bit data
 - ◆ 256-bit key schedule initialised from 128, 192, 256-bit keys
 - ◆ 16 round data computation using a complex highly non-linear function
 - ◆ two layers of designed S-P per round
 - ◆ same function also used in key schedule



LOKI97 Overview





LOKI97 Main Details

◆ data computation

$$R_i = L_{i-1} \text{ xor } f(R_{i-1} + SK_{3i-2}, SK_{3i-1})$$

$$L_i = R_{i-1} + SK_{3i-2} + SK_{3i}$$

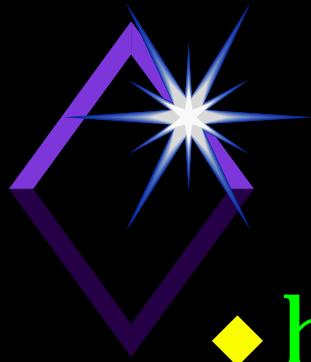
◆ key schedule

$$SK_i = K_{4i-1} \text{ xor}$$

$$f(K_{1i-1} + K_{3i-1} + nD, K_{2i-1})$$

$$K_{4i} = K_{3i-1}, K_{3i} = K_{2i-1}, K_{2i} = K_{1i-1}, K_{1i} = SK_i$$

$$D = \text{floor}((\text{sqrt}(5) - 1) \cdot 2^{63})$$

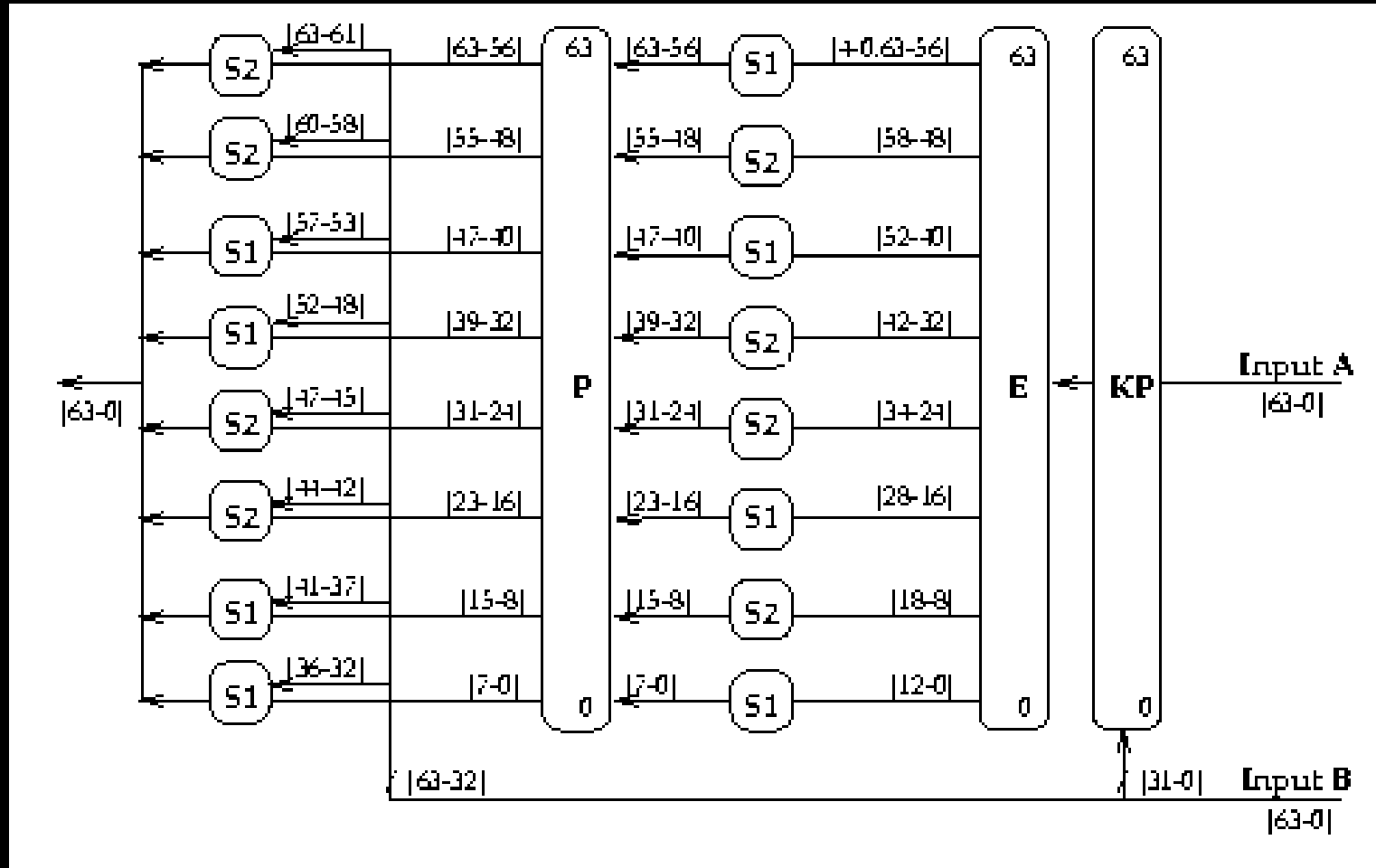


LOKI97 Round Function

- ◆ highly non-linear 64-bit function
 $f(A,B) = S_b(P(S_a(E(KP(A,B))))),B)$
- ◆ 2 columns each of 2 S-boxes
 - ◆ $S_a = [S1,S2,S1,S2,S2,S1,S2,S1]$
 - ◆ $S_b = [S2,S2,S1,S1,S2,S2,S1,S1]$
- ◆ regular perm P diffuses S_a outputs
 - ◆ fans S-box out to cover all S_b inputs
- ◆ keyed permutation KP to exchange selected pairs of bits set by input B



LOKI97 Function $f(A,B)$





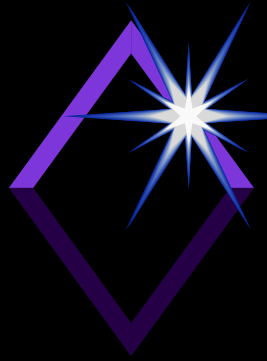
Rationale for S-boxes

- ◆ Want S-boxes which must
 - ◆ be balanced
 - ◆ be highly non-linear
 - ◆ satisfy strict avalanche criteria (SAC)
 - ◆ have good XOR profile



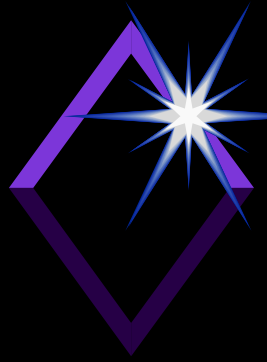
Rationale for S-boxes

- ◆ Cubing in odd Galois fields is proven to be good
 - ◆ $S(x) = x^3 \bmod p$, p irreducible polynomial in $GF2^n$
 - ◆ used $n = 11$ and 13
 - ◆ inverted input, truncated result to 8



Preliminary Analysis: Key Schedule

- ◆ no general linear relations for keys
 - ◆ no keys with repeated subkeys, eg 0
 - ◆ the first 4 subkeys can be coerced to 0 by solving suitable key schedule equations



Preliminary Analysis: Key Schedule

- ◆ the first 4 subkeys can be coerced to 0 by solving the key schedule equations

- ◆ SK_{1,2,3}=0,

$$K = [f(f(f(3D,0)+2D,0)+D, f(3D,0)) \mid f(f(3D,0)+2D,0) \mid f(3D,0) \mid 0]$$

- ◆ SK₁=0, $K = [f(D,0) \mid 0 \mid 0 \mid 0]$

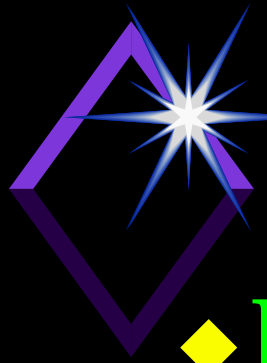
- ◆ SK_{1,2}=0,

$$K = [f(f(2D,0)+D,0) \mid f(2D,0) \mid 0 \mid 0]$$



Preliminary Analysis: S-Boxes

- ◆ XOR profiles very flat
 - ◆ S1 peak 64 (of 8192), zero peak 32
 - ◆ S2 peak Avalanche
 - ◆ 32 (of 2048), zero peak 16
 - ◆ ie standard 2 round characteristic has maximum $\text{Pr}(1/256)$
 - ◆ exhaustively tested all 1 bits changes
 - ◆ small number of 0 or 1 bit changes
 - ◆ similar to LOKI91



Preliminary Analysis: DC and LC

- ◆ Rijmen/Knudsen have suggested
 - ◆ DC attack using 2R characteristic via MSB
 - ◆ since XOR unchanged under addition
 - ◆ LC attack using a bias in f in directly keyed S_b layer



Possible Changes

- ◆ shift E down 2 bits, so MSB (63) is duplicated
 - ◆ should thwart DC attack
- ◆ consider alternatives for keying S_b to remove bias



Conclusions

- ◆ overview of LOKI97 design
 - ◆ previous work
 - ◆ rationale
 - ◆ goals
 - ◆ initial analysis
 - ◆ description
 - ◆ suggested alterations



Further Information

- ◆ on AES

<http://www.nist.gov/aes/>

- ◆ on LOKI97

<http://www.adfa.edu.au/~lpb/research/loki97/>