

An Introduction to Category Theory

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Categories

Definition (Category)

A category \mathbf{A} consists of

- a collection $ob(\mathbf{A})$ of **objects**
- for each $a, b \in ob(\mathbf{A})$ a collection $Hom_{\mathbf{A}}(a, b)$ of morphisms
- for each $a \in ob(\mathbf{A})$ a morphism $1_a \in Hom_{\mathbf{A}}(a, a)$ called the **identity** on a ($id_a = 1_a$)
- for each $a, b, c \in ob(\mathbf{A})$ a map

$$\begin{aligned} \circ : Hom_{\mathbf{A}}(a, b) \times Hom_{\mathbf{A}}(b, c) &\rightarrow Hom_{\mathbf{A}}(a, c) \\ (f, g) &\mapsto g \circ f \end{aligned}$$

Categories

Definition (Category)

- **associativity:** for each

$$f \in \text{Hom}_{\mathbf{A}}(a, b), g \in \text{Hom}_{\mathbf{A}}(b, c), h \in \text{Hom}_{\mathbf{A}}(c, d)$$

$$(h \circ g) \circ f = h \circ (g \circ f)$$

- **identity:** for each $f \in \text{Hom}_{\mathbf{A}}(a, b)$ we have

$$f \circ 1_a = f = 1_b \circ f$$

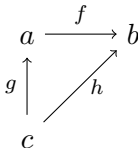


Examples

- **Set**
- **Grp**
- **Vect_k**
- **C with**
 - $ob(\mathbf{C}) = \{X\}$
 - $Hom_{\mathbf{C}}(X, X) := \{f : X \rightarrow X \mid f \text{ bijective}\}$
 - $(Hom_{\mathbf{C}}(X, X); \circ) \in ob(\mathbf{Grp})$

Notation of categories

Categories are noted as diagrams

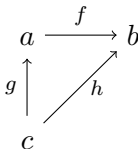


which show the objects and the morphisms between them.
For an morphism $f : a \rightarrow b$ we call

- $a = \text{dom}(f)$ the **domain** of f and
- $b = \text{cod}(f)$ the **codomain** of f .

Notation of categories

Categories are noted as diagrams



which show the objects and the morphisms between them.
We say the diagram **commutes** if

$$h = f \circ g.$$

Notation of categories

Categories are noted as diagrams

$$\begin{array}{ccc} a & \xrightarrow{f} & b \\ h \uparrow & & \uparrow j \\ c & \xrightarrow{g} & d \end{array}$$

which show the objects and the morphisms between them.
We say the diagram **commutes** if

$$f \circ h = j \circ g.$$



Properties of Categories

Definition

Let \mathbf{A} be a category. Then we call \mathbf{A}

- **small** if $ob(\mathbf{A})$ and $Hom_{\mathbf{A}}(a, b)$, for all $a, b \in ob(\mathbf{A})$, are actually sets
- **locally small** if $Hom_{\mathbf{A}}(a, b)$ is a set for each $a, b \in ob(\mathbf{A})$
- **large** otherwise

Recap group isomorphisms

Definition (Group homomorphisms)

Let $(G; \star)$ and $(H; \square)$ be groups. A group homomorphism is a map

$$\varphi : G \rightarrow H$$

s.t. for all $g_1, g_2 \in G$

$$\varphi(g_1 \star g_2) = \varphi(g_1) \square \varphi(g_2).$$

Recap group isomorphism

Definition (Group isomorphism)

Let $(G; \star)$ and $(H; \square)$ be groups. A group isomorphism is a map

$$\varphi : G \rightarrow H$$

s.t.

- φ is a group homomorphism and
- there exists a map $\varphi^{-1} : H \rightarrow G$ with

$$\varphi \circ \varphi^{-1} = id_H \text{ and } \varphi^{-1} \circ \varphi = id_G$$

Recap group homomorphisms

Example

Consider the groups

\star	0	1	2	3
0	0	1	2	3
1	1	2	3	0
2	2	3	0	1
3	3	0	1	2

and

\square	α	β	γ	δ
α	α	β	γ	δ
β	β	δ	α	γ
γ	γ	α	δ	β
δ	δ	γ	β	α

with

$$\varphi(0) = \alpha, \quad \varphi(1) = \beta, \quad \varphi(2) = \delta, \quad \varphi(3) = \gamma$$

the groups are isomorphic and φ is the isomorphism.



Product categories

Definition (Product category)

Let \mathbf{A} and \mathbf{X} be categories. Then the **product category** $\mathbf{A} \times \mathbf{X}$ of \mathbf{A} and \mathbf{X} looks like this:

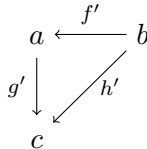
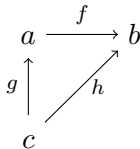
- $ob(\mathbf{A} \times_c \mathbf{X}) := ob(\mathbf{A}) \times ob(\mathbf{X})$
- $Hom_{\mathbf{A} \times_c \mathbf{X}}((a, x), (a', x')) := Hom_{\mathbf{A}}(a, a') \times Hom_{\mathbf{X}}(x, x')$
for $a, a' \in ob(\mathbf{A}), x, x' \in ob(\mathbf{X})$

Duality

Definition (Dual Category)

For each category \mathbf{A} there is a **dual category** \mathbf{A}^{op} s.t.

- $ob(\mathbf{A}^{op}) = ob(\mathbf{A})$
- for $f \in Hom_{\mathbf{A}}(a, b)$ there is $f' \in Hom_{\mathbf{A}^{op}}(b, a)$



Products

Definition (Product)

Let \mathbf{C} be a category and $X, Y \in ob(\mathbf{C})$. A **product of X and Y** is

$$(P, p_1 : P \rightarrow X, p_2 : P \rightarrow Y)$$

with $P \in ob(\mathbf{C})$ s.t. for all

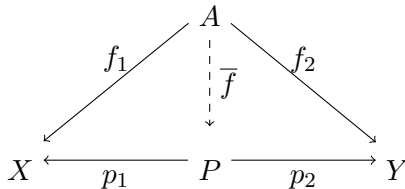
$$(A, f_1 : A \rightarrow X, f_2 : A \rightarrow Y)$$

with $A \in ob(\mathbf{C})$ there is a unique $\bar{f} : A \rightarrow P$ with

$$f_1 = p_1 \circ \bar{f}, \quad f_2 = p_2 \circ \bar{f}.$$

Products

Definition



commutes.

Coproducts

Definition (Coproduct)

Let \mathbf{C} be a category and $X, Y \in ob(\mathbf{C})$. A **coproduct of X and Y** is

$$(P, P \leftarrow X : i_1, P \leftarrow Y : i_2)$$

with $P \in ob(\mathbf{C})$ s.t. for all

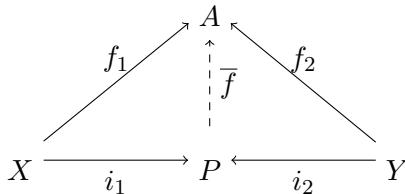
$$(A, A \leftarrow X : f_1, A \leftarrow Y : f_2)$$

with $A \in ob(\mathbf{C})$ there is a unique $A \leftarrow P : \bar{f}$ with

$$f_1 = \bar{f} \circ i_1, \quad f_2 = \bar{f} \circ i_2.$$

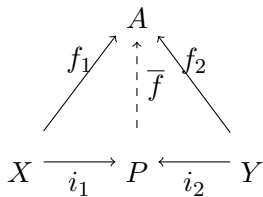
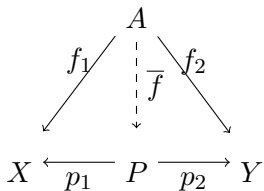
Coproducts

Definition



commutes.

Product/Coproduct



Functors

Definition (Functor)

Let \mathbf{A} and \mathbf{X} be categories. A Functor $T : \mathbf{A} \rightarrow \mathbf{X}$ is a map s.t.

- for all $a \in ob(\mathbf{A})$ we have $T(a) \in ob(\mathbf{X})$
- for all $a, a' \in ob(\mathbf{A})$ and $f \in Hom_{\mathbf{A}}(a, a')$ we have $T(f) \in Hom_{\mathbf{X}}(T(a), T(a'))$

which satisfies the following axioms:

- for each $a \in ob(\mathbf{A})$ we have $T(1_a) = 1_{T(a)}$
- for $f \in Hom_{\mathbf{A}}(a, a')$, $f' \in Hom_{\mathbf{A}}(a', a'')$ we have

$$T(f' \circ f) = T(f') \circ T(f)$$

Examples for Functors

- Identity functor: $id_{\mathbf{C}} : \mathbf{C} \rightarrow \mathbf{C}$
- Forgetful functor $U : \mathbf{Grp} \rightarrow \mathbf{Set}$ 'forgets group structure': A Group $(G; \star)$ is mapped to the underlying set G
- Free functor $F : \mathbf{Set} \rightarrow \mathbf{Grp}$ maps a set to the smallest group containing that set
- Bifunctors: $B : \mathbf{C}^{op} \times_{\mathbf{C}} \mathbf{C} \rightarrow \mathbf{D}$
 - Hom-bifunctor: $Hom : \mathbf{C}^{op} \times_{\mathbf{C}} \mathbf{C} \rightarrow \mathbf{Set}$
- Covariant Hom-functor: $Hom_{\mathbf{C}}(c, -) : \mathbf{C} \rightarrow \mathbf{Set}$
- Contravariant Hom-functor: $Hom_{\mathbf{C}}(-, c) : \mathbf{C}^{op} \rightarrow \mathbf{Set}$

Cones

Definition (Cone)

Let \mathbf{A} and \mathbf{I} be categories and $D : \mathbf{I} \rightarrow \mathbf{A}$ be a functor. A **cone on D** is an object $A \in \mathbf{A}$ and a family of morphisms $(A \xrightarrow{f_I} D(I))_{I \in \text{ob}(\mathbf{I})}$ s.t. for all $u : I \rightarrow J$ in \mathbf{I}

$$\begin{array}{ccc} & D(I) & \\ f_I \nearrow & & \downarrow Du \\ A & & \\ f_J \searrow & & \\ & D(J) & \end{array}$$

commutes.

Natural Transformations

Definition (Component)

Let \mathbf{A} and \mathbf{X} be categories, $F, G : \mathbf{A} \rightarrow \mathbf{X}$ be functors and $f \in \text{Hom}_{\mathbf{A}}(c, c')$ for $c, c' \in \text{ob}(\mathbf{A})$. A **component** is a morphism $\alpha_c \in \text{Hom}_{\mathbf{X}}(Fc, Gc)$ s.t.

$$\begin{array}{ccc} Fc & \xrightarrow{\alpha_c} & Gc \\ Ff \downarrow & & \downarrow Gf \\ Fc' & \xrightarrow{\alpha_{c'}} & Gc' \end{array}$$

commutes.

Natural Transformations

Definition (Natural Transformation)

Let \mathbf{A} and \mathbf{X} be categories and $F, G : \mathbf{A} \rightarrow \mathbf{X}$ be functors. A **natural transformation** $\alpha : F \Rightarrow G$ is a map assigning to every $c \in \text{ob}(\mathbf{A})$ its component $\alpha_c \in \text{Hom}_{\mathbf{X}}(Fc, Gc)$.



Adjoints

Definition (Adjoint)

Let \mathbf{A} and \mathbf{X} be categories and $F : \mathbf{A} \rightleftarrows \mathbf{X} : G$ be functors. Then F is **left adjoint** to G and G is **right adjoint** to F if

$$\varphi_{x,a} : \text{Hom}_{\mathbf{X}}(Fa, x) \cong \text{Hom}_{\mathbf{A}}(a, Gx)$$

and for

$$a, a' \in \text{ob}(\mathbf{A})$$

$$k \in \text{Hom}_{\mathbf{A}}(a', a)$$

$$x, x' \in \text{ob}(\mathbf{X})$$

$$h \in \text{Hom}_{\mathbf{X}}(x, x')$$

...

Adjoint

Definition (Adjoint)

...

$$\begin{array}{ccc}
 \text{Hom}_{\mathbf{X}}(Fa, x) & \xrightarrow{\varphi_{x,a}} & \text{Hom}_{\mathbf{A}}(a, Gx) \\
 \text{Hom}_{\mathbf{X}}(Fa, h) \downarrow & & \downarrow \text{Hom}_{\mathbf{A}}(a, Gh) \\
 \text{Hom}_{\mathbf{X}}(Fa, x') & \xrightarrow{\varphi_{x',a}} & \text{Hom}_{\mathbf{A}}(a, Gx')
 \end{array}$$

and

$$\begin{array}{ccc}
 \text{Hom}_{\mathbf{X}}(Fa, x) & \xrightarrow{\varphi_{x,a}} & \text{Hom}_{\mathbf{A}}(a, Gx) \\
 \text{Hom}_{\mathbf{X}}(Fk, x) \downarrow & & \downarrow \text{Hom}_{\mathbf{A}}(k, Gx) \\
 \text{Hom}_{\mathbf{X}}(Fa', x) & \xrightarrow{\varphi_{x,a'}} & \text{Hom}_{\mathbf{A}}(a', Gx)
 \end{array}$$

commute.

Example

$$V : \mathbf{Set} \rightleftarrows \mathbf{Vect}_k : U$$

take $X \in \text{ob}(\mathbf{Set})$, $W \in \text{ob}(\mathbf{Vect}_k)$

$$\varphi_{X,W} : \text{Hom}_{\mathbf{Vect}_k}(V(X), W) \cong \text{Hom}_{\mathbf{Set}}(X, U(W))$$



(Co)units

Definition ((Co)unit)

Every adjunction $F : \mathbf{A} \rightleftarrows \mathbf{X} : G$ yields two natural transformations:

- the **unit** natural transformation

$$\eta : 1_{\mathbf{A}} \Rightarrow G \circ F$$

$$\eta_A : A \rightarrow GFA$$

- the **counit** natural transformation

$$\epsilon : F \circ G \Rightarrow 1_{\mathbf{X}}$$

$$\epsilon_X : FGX \rightarrow X$$



Literature

- (*) Basic Category Theory - Tom Leinster
- (*) Categories and Computer Science - RFC Walters
- (***) Categories for the working mathematician - Saunders Mac Lane
- (***) ncatlab.org
- (**) en.wikipedia.org